

# Lesson Plan: Complexity Olympics

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**Year Group:** 9 | **Duration:** 50 minutes | **Topic:** Algorithmic Complexity / Big-O Notation

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## 1. Overview

**Core Concept:** Algorithmic complexity — how the number of steps grows as the input size grows.  $O(1)$ ,  $O(\log n)$ ,  $O(n)$ ,  $O(n^2)$  demonstrated through physical activities.

### Learning Objectives:

- Experience  $O(1)$ ,  $O(\log n)$ ,  $O(n)$ , and  $O(n^2)$  through physical activities
- Plot growth curves and identify their shapes (flat, logarithmic, linear, quadratic)
- Connect complexity classes to real algorithm choices
- Understand why  $O(n^2)$  algorithms become unusable at large scale

### Key Vocabulary:

Term	Definition
Complexity	How the number of steps grows as input size grows
Big-O notation	A way of expressing the growth rate of an algorithm
$O(1)$	Constant time — always the same number of steps
$O(\log n)$	Logarithmic time — steps grow very slowly
$O(n)$	Linear time — steps grow proportionally with input
$O(n^2)$	Quadratic time — steps grow as the square of input size
Input size ( $n$ )	The amount of data being processed

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## 2. Before the Lesson

### Print:

- [resource-station-cards.md](#) — 1 per station (laminated if possible)
- [worksheet-results-recording.md](#) — 1 per student

### Room Setup:

- 4 stations around the room, each with its instruction card
  - Each station needs a small stack of numbered index cards
  - Groups of 5–6 students, rotating every ~7 minutes
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## 3. Timed Lesson Flow

## 0–5 min — Introduce the Olympics

1. Today: four events. Each measures how the WORKLOAD grows as the INPUT gets bigger.
2. At each station: read the card, do the activity for small  $n$ , medium  $n$ , large  $n$ . Record on worksheet.

## 5–35 min — Rotation Through Stations

Groups rotate, ~7 minutes per station. Teacher circulates.

**Station 1 —  $O(1)$ :** Direct lookup. Given a position number, retrieve that card instantly. Always 1 step, regardless of stack size.

**Station 2 —  $O(\log n)$ :** Binary search on a sorted stack. Count guesses for different stack sizes.

**Station 3 —  $O(n)$ :** Find the maximum in an unsorted pile. Must look at every card.

**Station 4 —  $O(n^2)$ :** Check every pair of cards for duplicates. Count total comparisons.

## 35–45 min — Graph

Students complete the worksheet graph — plotting all 4 curves on the same axes.

## 45–50 min — Debrief

- Which algorithm would you want for 1 million items?
- Why does  $O(n^2)$  become unusable at scale?
- *"If sorting 10 items takes 100 steps with  $O(n^2)$  sort, how many steps for 100 items? For 1000?"*

## 4. Teacher Facilitation Notes

### What to look for:

- Station 4: students undercount pairs — remind them: card 1 pairs with ALL others, then card 2 pairs with remaining others. Formula:  $n \times (n-1) \div 2$
- Station 1: students feel like they're "cheating" — that IS the point.  $O(1)$  means the data structure gives instant access (like an array index or dictionary key)
- Help students predict before doing: "If  $n=5$  takes 10 comparisons, how many will  $n=10$  take?" The doubling to 40 (not 20) surprises people

### Common misconceptions:

- $O(1)$  means fast — no, it means constant. An  $O(1)$  operation could still take 1 second.  $O(n)$  on 3 items might be faster.
- Big- $O$  is the exact number of steps — no, it's the growth RATE (proportionality). Constants and lower-order terms are ignored.

## 5. Extension Tasks

1. For  $n=100$ , station 4 takes 4,950 comparisons. For  $n=1000$ : ~500,000. Plot these. Why is  $O(n^2)$  impractical at scale?

2. If a computer does 1 billion operations per second, how long does  $O(n^2)$  take for  $n=1,000,000$ ?
  3. Research: what is  $O(n \log n)$ ? Which sorting algorithms achieve this?
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## 6. Key Takeaway

**The growth rate of an algorithm matters as much as its correctness.  $O(n^2)$  algorithms become impractical at scale.  $O(\log n)$  algorithms scale magnificently. Choosing the right complexity class is a core engineering decision.**